

CSC263 Tutorial #11

MSTs

March 31, 2023

Things covered in this tutorial

- ★ What's the **Disjoint Set** ADT?
- ★ How is it implemented using trees?

Disjoint Sets

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Disjoint Sets

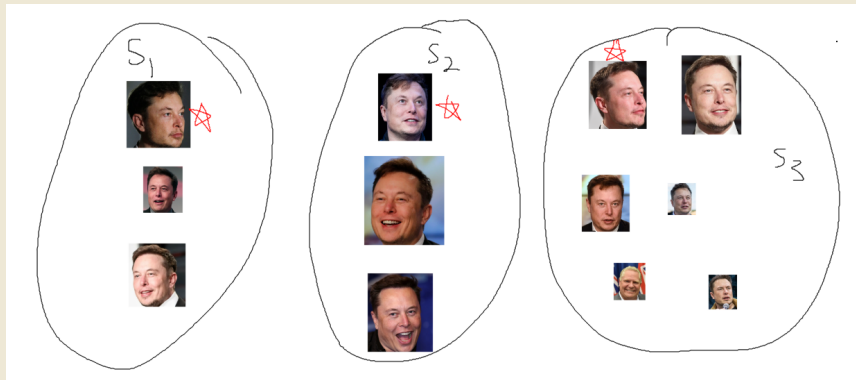
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Answer:

- ★ **MakeSet**(x): Add a singleton set $\{x\}$ to the collection of disjoint sets.
- ★ **FindSet**(x): Return the representative of the set containing x .
- ★ **Union**(x, y): Combine the two disjoint sets S_i containing x and S_j containing y respectively (and ignore if $S_i = S_j$).

Disjoint Sets



Disjoint Sets

Task: Apply the following disjoint set operations, and draw the result as shown in the previous slide.

★ MakeSet(Chungi)

★ MakeSet(Winni)

★ MakeSet(Bob)

★ MakeSet(727)

★ MakeSet(Elong)

★ MakeSet(420)

★ MakeSet(Mogus)

★ Union(420, Mogus)

★ Union(Mogus, Bob)

★ Union(Winni, 727)

★ Union(Winni, Bob)

★ Union(Winni, Chungi)

Disjoint Sets with Trees!

Recall how we tried to implement disjoint sets with linked lists, but that wasn't very efficient...

Linked list with **pointer to head**

MakeSet and **FindSet** are fast, **Union** now becomes the time-consuming one, especially if appending a long list.

Amortized analysis: The total cost of a sequence of m operations.

- Bad sequence: $m/2$ **MakeSet**, then $m/2 - 1$ **Union**, then **1** **FindSet**.
 - ◆ Always let the longer list append, like 1 append 1, 2 append 1, 3 append 1, ..., $m/2 - 1$ append 1.
- Total cost: $\Theta(1+2+3+\dots+m/2 - 1) = \Theta(m^2)$

Linked list with **union-by-weight**

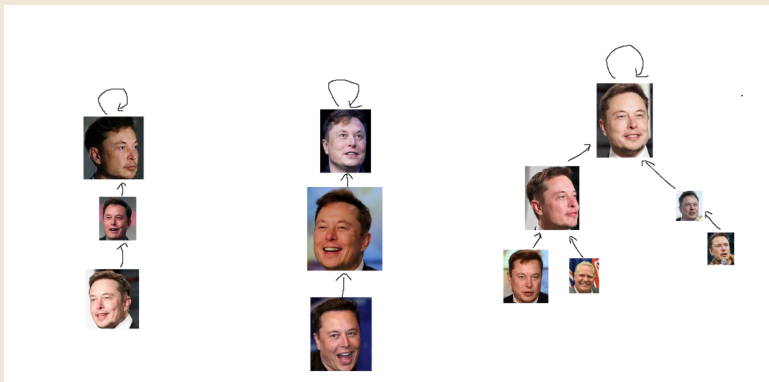
Union-by-weight sounds like a simple heuristic, but it actually provides significant improvement.

For a sequence of m operations which includes n **MakeSet** operations, i.e., n elements in total, the total cost is $\mathbf{O(m + n \log n)}$

i.e., for the previous sequence with $m/2$ **MakeSet** and $m/2 - 1$ **Union**, the total cost would be $\mathbf{O(m \log m)}$, as opposed to $\mathbf{\Theta(m^2)}$ when without union-by-weight.

Disjoint Sets with Trees!

We implement Disjoint Set using an “Inverted Tree”:



We use two tricks to speed up the operations: **Union By Rank** and **Path Compression**.

Disjoint Sets with Trees!

Task: Using an inverted tree implementation with Union by Rank and Path Compression, apply the following disjoint set operations.

★ MakeSet (C)

★ MakeSet (W)

★ MakeSet (B)

★ MakeSet (7)

★ MakeSet (E)

★ MakeSet (4)

★ MakeSet (M)

★ Union(4, M)

★ Union(M, B)

★ Union(W, 7)

★ Union(7, M)

★ Findset(7)

★ Findset(M)

★ Union(W, C)

Disjoint Sets with Trees!

Task: Complete the tutorial activity! (Why can't we just traverse an inverted tree?)

Disjoint Sets with Trees!

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Hint: The two attributes needed are named `node.next` and `node.tail` respectively.

Goodbye!



Good luck on exams!