# CSC263 Tutorial #4 AVL Trees

February 3, 2023

## Things covered in this tutorial

- \* What's an AVL tree?
- $\star\,$  How do I insert into or delete from an AVL tree?
- \* How do I rotate?
- $\star$  How can I make the AVL tree very slow?<sup>1</sup>

<sup>&</sup>lt;sup>1</sup>Still  $\mathcal{O}(\log n)$  though...

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To delete a node from a BST, there are three cases:

- \* The node has no children: delete the node.
- \* The node has one child: delete the node, promote the child.
- \* The node has two children: swap the node with its successor, then delete the node (which has 0 or 1 children).

Height of a binary tree: starts from 0.



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Insert(root, key, value):
    # BST insert as usual
    BST_insert(root, key, value)

    # Fix balance, update height
    root.balance_factor = root.right.height - root.left.height
    if root.balance_factor < -1 or root.balance_factor > 1:
        fix_imbalance(D) # Rotations!
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- \* Insert(115).



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Perform the following:

- \* Insert(85).
- $\star$  Insert(115).

For Insert(115), we needed to do a left-right rotation.













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